## Programming 1C: A branch-and-bound approach to prune the search space

## 1 The mixture sequence problem

A sequence of *n* natural numbers  $C = \langle c_1, c_2, \ldots, c_n \rangle$  is said to be a *mixture* of a sequence of *n* natural numbers  $A = \langle a_1, a_2, \ldots, a_n \rangle$  and another sequence of *n* natural numbers  $B = \langle b_1, b_2, \ldots, b_n \rangle$  if and only if for every  $1 \le i \le n$  we have  $c_i = a_i$  or  $c_i = b_i$ . In other words, for *i* from 1 to *n*, we pick either  $a_i$  or  $b_i$  and make it  $c_i$ . We say  $\sum_{1 \le i \le n} c_i$  is a *mixture sum* of *A* and *B*.

The mixture sequence problem is defined as follows:

- Given two sequences of n natural numbers of  $A = \langle a_1, a_2, \ldots, a_n \rangle$  and  $B = \langle b_1, b_2, \ldots, b_n \rangle$ ,
- and given a natural number m,
- determine whether there exists a mixture C of A and B such that the the mixture sum of C equals m, and
- if such a mixture sequence does exist print out the mixture sequence; otherwise print out a message saying there is no such mixture sequence.

## 2 A Branch-and-Bound Approach

The following is a possible branch-an-bound approach to solve the problem:

- We refer to an empty sequence  $C' = \langle \rangle$  of no elements as a partial mixture sequence of length 0. For  $1 \leq l < n$ , let's refer to a sequence  $C' = \langle c_1, c_2, \ldots, c_l \rangle$  as a partial mixture sequence (of A and B) of length l if and only if  $c_i = a_i$  or  $c_i = b_i$  for  $1 \leq i \leq l$ . We use the notation PartialSum(C')to refer to the sum of all the elements in the partial mixture sequence C'.
- For  $0 \leq l \leq n$ , let's define  $L_i = \sum_{l+1 \leq k \leq n} \min(a_k, b_k)$  and  $U_l = \sum_{l+1 \leq k \leq n} \max(a_k, b_k)$ . In other words,  $L_l$  is the minimum (lower bound) of the sum of the l+1 th element to the *n* th (the last) element in any mixture sequence *C* derived from *A* and *B* while  $U_l$  is the maximum (upper bound) of the sum of the l+1 th element to the *n*th (the last) element in any mixture sequence *C* derived from *A* and *B* while  $U_l$  is the maximum (upper bound) of the sum of the l+1 th element to the *n*th (the last) element in any mixture sequence *C* derived from *A* and *B*.
- Observation 1: We can search for a mixture sequence C with a target mixture sum m by starting from an empty partial mixture sequence C' of length l = 0 and then iteratively trying to extend the current partial mixture sequence C' = ⟨..., c<sub>l</sub>⟩ of length l one step further by appending either a<sub>l+1</sub> or b<sub>l+1</sub> to the end of the current partial mixture sequence C'. Naively, for a partial mixture sequence C' of length l there are two possible branches

to explore for further extensions into a final mixture sequence C with the target mixture sum m. However, Observation 2 below shows that at time we can completely cut off any search along these two branches completely by checking the bounds  $L_l$  and  $U_l$  against PartialSum(C') and the target sum m.

• Observation 2: Note that whatever way we extend a partial mixture sequence  $C' = \langle \ldots, c_l \rangle$  into a mixture sequence  $C = \langle c_1, c_2, \ldots, c_n \rangle$ , the resulting mixture sum must be at least  $PartialSum(C') + L_l$  and at most  $PartialSum(C') + U_l$ . If the target sum *m* does not fall into the range of  $[PartialSum(C') + L_l, PartialSum(C') + U_l]$ , we know that it is impossible to extend the current partial mixture sequence C' into a mixture sequence ending in the target mixture sum of *m*. In this case, we can completely cut off any further search effort to extend the current partial mixture sequence C'.

## 3. Implementation and experiments to do:

Implement the branch-and-bound scheme above for solving the mixture sequence problem. **Note**: You may consider using <u>this</u> recursive exhaustive branching search implementation as the basis and then incorporate the branch-and-bound scheme into it.

**Experiment #A**: Run your new program over at least the test cases 1A~1D in Programming #1B. Please report whether your new program can get the same results correctly as your program for Programming #1B did.

**Experiment #B**: Consider the following two sequences:

A =<18, 29, 31, 42, 55, 66, 71, 85, 91, 103, 114, 128, 135, 140, 155, 81, 93, 17, 24, 65, 39, 27, 58, 19, 130, 141, 182, 153, 104, 115, 18, 29, 31, 42, 55, 66, 71, 85, 91, 103, 114, 128, 135, 140, 155> and

B =<81, 93, 17, 24, 65, 39, 27, 58, 19, 130, 141, 182, 153, 104, 115, 18, 29, 31, 42, 55, 66, 71, 85, 91, 103, 114, 128, 135, 140, 155, 81, 93, 17, 24, 65, 39, 27, 58, 19, 130, 141, 182, 153, 104, 115 >.

Use your new program to search for a mixture sequence with the sum of (i) 655, (ii) 3869, (iii) 3901, (iv) 4139, (v) 99991 respectively.

Report your findings on cases that the program can finish within 1 day. For each case, please report whether your new program can find a mixture sequence (and include the sequence if it does find one). How much time on average does it take to finish one case?

**Experiment #C**: Try to use your old program for Programming #1B to repeat what you are required to do in the Experiment B above. **Report your findings on cases that the program can finish within 1 day.**