

## Test 2: Divide and Conquer

For #1, #2(i), #3(i), #4, and #5 below, you need to describe (i) the divide-and-conquer algorithm you propose for solving the problem, (ii) the resulting recurrence relation in the form of  $T(n) = a \cdot T(n/b) + n^c$  regarding the time complexity  $T(n)$  and (iii) the time complexity  $T(n)$  in closed form according to the master method in the handout and the actual coefficients  $a$ ,  $b$  and  $c$ , in the recurrence relation in (ii).

1. Given an array  $A$  of  $n$  elements  $A[1], A[2], \dots, A[n]$  and you are also told the values of these  $n$  elements have the following V shape property: there is a unique (untold) index  $k$  in the range of  $[1, n]$  such that  $A[i] > A[i+1]$  for every index  $i$  in the range of  $[1, k-1]$  and  $A[i] < A[i+1]$  for every index  $i$  in the range of  $[k, n-1]$ . Obviously  $A[k]$  is the minimum among all the values in the entire array. **Devise a good divide-and-conquer algorithm to find this minimal value. The time complexity of the algorithm should be better than  $\Theta(n)$ . 10 points.**

2.

(i) Given an array  $A$  of  $n$  elements  $A[1], A[2], \dots, A[n]$  storing arbitrarily large integers, the maximum value gap in this array is the maximum of  $|A[i] - A[j]|$  over all possible index pairs  $i$  and  $j$  **where  $i > j$** . **Devise a good divide-and-conquer algorithm to find the maximum value gap. Your algorithm should be more efficient than an  $\Theta(n^2)$  brute force search to find the maximum value gap. 10 points.**

(ii) Implement your algorithm as a running program that can ask the user to enter the  $n$  values and then determine the maximum value gap. **5 points.**

3.

(i) Given an array  $A$  of  $n$  elements  $A[1], A[2], \dots, A[n]$  storing arbitrarily large integers, we say a pair of values of  $A[i]$  and  $A[j]$  is a significant inversion if and only if  $i < j$  and  $A[i] > 2 \cdot A[j]$ . **Devise a good divide-and-conquer to find the total number of significant inversion pairs in  $A$ . Your algorithm should be more efficient than an  $\Theta(n^2)$  brute force search that examine all the  $\Theta(n^2)$  pairs to find the number of significant inversions. 10 points.**

(ii) Implement your algorithm as a running program that can ask the user to enter the  $n$  values and then determine the number of significant inversion pairs. **5 points.**

4. Consider a two dimensional array  $A$  of  $n$  elements  $A[i][j]$  where  $i$  and  $j$  are indices in the range of  $[1, n^{1/2}]$ . All the elements  $A[i][j]$  are **distinct** integers. Two elements  $A[i][j]$  and  $A[k][l]$  are neighbors if and only if  $|i-k| + |j-l|$  equals 1. An element  $A[i][j]$  is a local minimum if and only if the integer stored in  $A[i][j]$  is smaller than all of its neighbors. **Devise a good divide-and-conquer to find a local minimum. Your algorithm should be more efficient than an  $\Theta(n)$  brute force search that examine all the  $n$  elements and their neighbors. 10 points.**

**5. Problem for bonus points.** Given an array  $A$  of  $n$  elements  $A[1], A[2], \dots, A[n]$  storing arbitrarily large integers, we say an integer  $k$  is the majority value in this array if and only if more than  $n/2$  of the elements in  $A$  store this particular value  $k$ . Note that there may or may not be a majority value. Assuming that checking whether two integers are equal takes only constant time and  $n$  is a power of 2. **Devise a good divide-and-conquer algorithm that uses no more than  $2n$  comparisons to determine whether the majority value does exist in the array and if it does exist find out this majority value.**  
**Bonus 10 points.**